

**Discrete Graphical Models  
— An Optimization  
Perspective**



# Discrete Graphical Models — An Optimization Perspective

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## Contents

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<b>1</b>	<b>Introduction to Inference for Graphical Models</b>	<b>7</b>
1.1	Basic definitions . . . . .	9
1.2	Probabilistic interpretation . . . . .	12
1.3	Combinatorial complexity of MAP-inference . . . . .	13
1.4	Bibliography and further reading . . . . .	15
<b>2</b>	<b>Acyclic Graphical Models</b>	<b>17</b>
2.1	MAP-Inference with dynamic programming . . . . .	18
2.2	Computation of min-marginals . . . . .	22
2.3	On dynamic programming for cyclic graphs . . . . .	24
2.4	Bibliography and further reading . . . . .	25
<b>3</b>	<b>Background: (Integer) Linear Programs and Their Geometry</b>	<b>26</b>
3.1	Optimization problems . . . . .	26
3.2	Linear constraints and polyhedra . . . . .	30
3.3	Linear programs . . . . .	39
3.4	Integer linear programs . . . . .	41
3.5	Linear program relaxation . . . . .	45
3.6	Bibliography and further reading . . . . .	47

<b>4</b>	<b>Energy Minimization as Integer Linear Program</b>	<b>48</b>
4.1	MAP-inference as ILP . . . . .	48
4.2	ILP properties of the MAP-inference problem . . . . .	53
4.3	Properties of the local polytope relaxation . . . . .	54
4.4	Bibliography and further reading . . . . .	59
<b>5</b>	<b>Background: Basics of Convex Analysis</b>	<b>60</b>
5.1	Convex optimization problems . . . . .	60
5.2	Subgradient . . . . .	67
5.3	Lagrange duality . . . . .	71
5.4	Lagrange relaxation of integer linear programs . . . . .	78
5.5	Bibliography and further reading . . . . .	82
<b>6</b>	<b>Lagrange Duality for MAP-inference</b>	<b>84</b>
6.1	Reparametrization and Lagrange dual . . . . .	85
6.2	Primal solutions from the dual problem . . . . .	92
6.3	Dual optimality for acyclic graphical models . . . . .	106
6.4	Bibliography and further reading . . . . .	108
<b>7</b>	<b>Background: Basics of Non-Smooth Convex Optimization</b>	<b>109</b>
7.1	Gradient descent . . . . .	109
7.2	Sub-gradient method . . . . .	113
7.3	Coordinate descent . . . . .	116
7.4	Bibliography and further reading . . . . .	119
<b>8</b>	<b>Subgradient and Coordinate Descent for MAP-Inference</b>	<b>120</b>
8.1	Primal coordinate descent . . . . .	121
8.2	Dual sub-gradient method . . . . .	125
8.3	Min-sum diffusion as block-coordinate ascent . . . . .	128
8.4	Convergence of dual block-coordinate ascent methods . . . . .	139
8.5	Abstract convergence theorem . . . . .	139
8.6	Convergence of diffusion algorithms . . . . .	142
8.7	Empirical comparison of algorithms . . . . .	148
8.8	Bibliography and further reading . . . . .	152



<b>9 Lagrange (Dual) Decomposition</b>	<b>156</b>
9.1 Lagrange decomposition in a nutshell . . . . .	157
9.2 Lagrange decomposition for grid graphs . . . . .	158
9.3 General graph decomposition scheme . . . . .	162
9.4 Equivalence of all acyclic decompositions . . . . .	174
9.5 Bibliography and further reading . . . . .	180
<b>10 Maximization of the Decomposition-Based Dual</b>	<b>182</b>
10.1 Subgradient method . . . . .	183
10.2 Tree Reweighted Message Passing (TRW-S) . . . . .	186
10.3 Empirical comparison of algorithms . . . . .	200
10.4 Bibliography and further reading . . . . .	203
<b>11 Min-Cut/Max-Flow Based Inference</b>	<b>209</b>
11.1 Submodular functions . . . . .	210
11.2 Submodular pairwise energies . . . . .	214
11.3 Binary submodular problems as min-cut . . . . .	224
11.4 Transforming multi-label problems into binary ones . . . . .	231
11.5 Move-making algorithms . . . . .	236
11.6 Bibliography and further reading . . . . .	242
<b>12 Relaxed Binary Energy as <i>st</i>-Min-Cut</b>	<b>243</b>
12.1 Half-integrality of local polytope . . . . .	244
12.2 LP relaxation as min-cut . . . . .	246
12.3 Persistency of the binary local polytope . . . . .	249
12.4 Bibliography and further reading . . . . .	254
<b>Acknowledgements</b>	<b>255</b>
<b>References</b>	<b>257</b>

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## ABSTRACT

This monograph is about combinatorial optimization. More precisely, about a special class of combinatorial problems known as *energy minimization* or *maximum a posteriori (MAP) inference in graphical models*, closely related to *weighted* and *valued constraint satisfaction problems* and having tight connections to *Markov random fields* and *quadratic pseudo-boolean optimization*. What distinguishes this monograph from a number of other monographs on graphical models is its focus: It considers graphical models, or, more precisely, MAP-inference for graphical models, purely as a combinatorial optimization problem. Modeling, applications, probabilistic interpretations and many other aspects are either ignored here or find their place in examples and remarks only.

Combinatorial optimization as a field is largely based on five fundamental topics: (i) integer linear programming and polyhedral optimization; (ii) totally unimodular matrices and the class of min-cost-flow problems; (iii) Lagrange decompositions and relaxations; (iv) dynamic programming and (v) submodularity, matroids and greedy algorithms. Each of these topics found its place in this monograph, although to a variable extent. The covering of each respective topic

reflects its importance for the considered MAP-inference problem.

Since optimization is the primary topic of this monograph, we mostly stick to the terminology widely used in optimization and where it was possible we tried to avoid the graphical models community-specific technical terms. The latter differ from sub-community to sub-community and, in our view, significantly complicate the information exchange between them.

The same holds also for the presentation of material in this monograph. If there is a choice when introducing mathematical constructs or proving statements, we prefer more general mathematical tools applicable in the whole field of operations research rather than to stick to graphical model-specific constructions. We additionally provide the graphical model-specific constructions if it turns out to be easier than the more general one. This way of presentation has two advantages. A reader familiar with a more general technique can grasp the new material faster. On the other hand, the monograph may serve as an introduction to combinatorial optimization for readers unfamiliar with this subject. To make the monograph even more suitable for both categories of readers we explicitly separate the mathematical optimization background chapters from those specific to graphical models.

We believe, therefore, that the monograph can be useful for undergraduate and graduate students studying optimization or graphical models, as well as for experts in optimization who want to have a look into graphical models. Moreover, we believe that even experts in graphical models can find new views on the known facts as well as a novel presentation of less known results in the monograph. These are for instance (i) a simple and general proof of equivalence of different acyclic Lagrange decompositions of a graphical model; (ii) a general scheme for analyzing convergence of

dual block-coordinate descent methods; (iii) a short and self-contained analysis of a linear programming relaxation for binary graphical models, its persistency properties and its relation to quadratic pseudo-boolean optimization.

The present monograph is based on lectures given to undergraduate students of Technical University of Dresden and Heidelberg University. The selection of material is done in a way that it may serve as a basis for a semester course.

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## Notation

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To simplify reading of the monograph, some frequently used notations are collected here. Some of them, which we assume to be quite standard, are used without additional notice in the text. Others, typically more specialized, are introduced in the monograph. For those we point out the section and the page they are defined in.

**Standard notation**

$\mathbb{N}$	the set of natural numbers
$\mathbb{Z}$	the set of integer numbers
$\mathbb{R}$	the set of real numbers
$\mathbb{R}^n$	an $n$ -dimensional vector space over the field of real numbers
$\mathbb{R}_+^n$	the set of vectors with non-negative coordinates in $\mathbb{R}^n$ , i.e. $\{x \in \mathbb{R}^n : x_i \geq 0, i = 1, 2, \dots, n\}$ ; for $n = 1$ the notation simplifies to $\mathbb{R}_+$
$x \in \mathcal{A}^{\mathcal{B}}$	For any set $A$ and any finite set $\mathcal{B}$ , this notation stands for a vector $x$ with $ \mathcal{B} $ coordinates indexed by elements of $\mathcal{B}$ , for each $b \in \mathcal{B}$ it holds that $x_b \in \mathcal{A}$ . The only exception from this rule is the notation $\Delta^{\mathcal{B}}$ , see below.
$x \geq y$	comparison operations are applied coordinate-wise to vectors and point-wise to functions
$\llbracket \cdot \rrbracket$	denotes the Iverson brackets, that is, for any predicate $A$ it holds that $\llbracket A \rrbracket = 1$ if $A$ is true, otherwise $\llbracket A \rrbracket = 0$
$\langle c, x \rangle$	the inner product, i.e. $\langle c, x \rangle = \sum_{i=1}^n c_i x_i$
$\nabla f$	gradient of the function $f$
$O(\cdot)$	for two functions $f: \mathbb{N} \rightarrow \mathbb{N}$ and $g: \mathbb{N} \rightarrow \mathbb{N}$ one writes $f = O(g)$ , if there is a constant $c > 0$ and a number $n_0 \in \mathbb{N}$ such that $f(n) \leq c \cdot g(n)$ for all $n \geq n_0$

**Standard abbreviations**

w.r.t.	with respect to
w.l.o.g.	without loss of generality
s.t.	subject to

**Notation defined in the monograph**

$\mathcal{N}_b(u)$	set of graph vertexes incident to vertex $u$ , see §1.1, page 9
$\delta_{\mathcal{G}}(y), \delta(y)$	binary representation of the labeling $y$ , i.e. a binary vector with non-zero coordinates corresponding to the labels $y_u, u \in \mathcal{Y}_u$ , and label pairs $(y_u, y_v), uv \in \mathcal{E}$ , see page 49;
$\mathcal{I}$	the set of indexes of the cost vector of a graphical model; $ \mathcal{I} $ is equal to the number of coordinates of the cost vector, see §1.1 on page 11
$\text{vrtx}(P)$	set of vertexes of the polyhedron $P$ , see Definition 3.19 on page 32
$\Delta^n$	$n$ -dimensional simplex, see Definition 3.21 on page 33
$\Delta^{\mathcal{X}}$	$ \mathcal{X} $ -dimensional simplex, with coordinates indexed by elements of $\mathcal{X}$ , see Definition 3.21 on page 33
$\text{conv}(X)$	convex hull of $X$ , see Definition 3.28 on page 34
$\text{mi}[\theta_w]$	binary vector with non-zero coordinates corresponding to locally minimal values of the cost vector $\theta_w$ , see page 95
$\text{nz}[\mu]$	binary vector with non-zero coordinates corresponding to the non-zero coordinates of $\mu$ , see page 95
$\text{cl}(\xi)$	arc-consistency closure of a binary vector $\xi$ , see Definition 6.11 on page 99
$\mathcal{J}$	the set of indexes of the Lagrange dual vector for the MAP-inference problem; $ \mathcal{J} $ is equal to the number of coordinates in the dual vector, see §6.1 on page 85
$\langle \frac{1}{2} \rangle, \langle 0.7 \rangle$	angular brackets are used in figures for coordinates of primal relaxed solutions, see e.g. Figure 4.1, 4.2, 6.4, 12.3

# 1

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## Introduction to Inference for Graphical Models

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There are many problems in computer science, which can be formulated in the form of so-called *Discrete Graphical Models*. Examples can be found in bio-informatics, communication theory, statistical physics, computer vision, signal processing, information retrieval and machine learning.

Discrete graphical models as a modeling tool naturally appear when

- the target object (the object we model) consists of *many small parts*,
- each part must be labeled by a label from a *finite set*, and
- parts (and, therefore, their labels) are *mutually dependent*.

**Example 1.1** (Image segmentation). Image segmentation is a classical image analysis task: Each pixel of an input image must be assigned a label of an object visible in the image. For instance, if we consider images of street scenes, these labels could belong to the set {pedestrian, car, tree, building}.

The target object is an image, i.e. a two-dimensional array of pixels. Each pixel constitutes an elementary part of the image and must be



labeled with a label from a finite set. The simplest assumption about image segments, i.e. groups of pixels having the same label, is the so called “compactness assumption”. It states that it is more probable that neighboring pixels are labeled with the same label than with different ones.

**Example 1.2** (Depth reconstruction). Depth reconstruction is another important image analysis problem. In the classical setting there are (at least) two images taken from different viewpoints. The task is to match pixels from these two images to each other. Assuming the positions of cameras and their focal lengths are known, this allows us to estimate depth of the scene, which was photographed with the cameras.

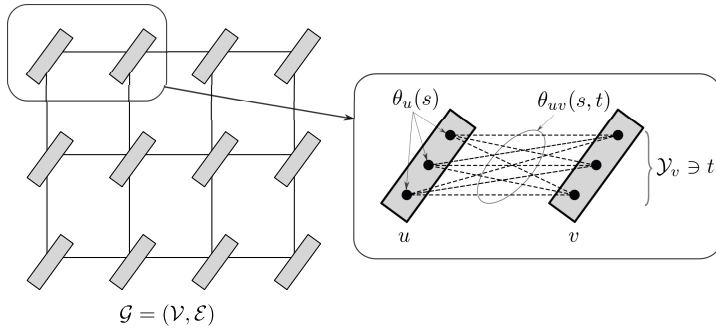
As in the previous example, the target object is a two-dimensional pixel array, where each pixel constitutes an elementary part of the object and must be labeled with a label from a finite set. Here, the meaning of the labels is different: Each label represents depth information of the associated pixel in an image, i.e. how far the depicted observation is placed from the camera. Usually the set of labels is chosen as natural numbers in a given interval, for instance,  $\{0, 1, \dots, 255\}$ .

Assuming that the observed surface is smooth, one would expect the difference  $|s - s'|$  between labels  $s$  and  $s'$  in neighboring pixels to be small. The opposite would mean a jump in depth, or, in other words, non-smoothness of the surface.

**Example 1.3** (Cell tracking problem in bio-imaging). Given is a sequence of images that show the development of a living organism from an early embryo consisting of only a few cells to a fully grown animal. During this sequence, the images show moving and splitting cells.

Under the assumption that the image is already *pre-segmented*, i.e. the cells were already found in each image, the task at hand is to track each individual cell and its descendants from the first to the last frame.

The cells are the elementary parts of the considered object. Each cell in a given image frame is labeled with pointers to one or two cells in the next frame. One pointer means that the cell only moved, and two pointers correspond to a cell division. The simplest tracking model forbids two different cells to have the same descendants. This rule defines dependencies between object parts.



**Figure 1.1:** Example of a graphical model with grid structure. On the left, graph nodes are denoted with inclined rectangles, lines connecting nodes correspond to graph edges. On the right, two neighboring nodes are shown. Black circles inside rectangles correspond to the labels  $s$  in the node  $u$  (left rectangle) and  $t$  in the node  $v$  (right rectangle). Dashed lines correspond to each label pair  $s, t$  with an assigned pairwise cost  $\theta_{uv}(s, t)$ .

## 1.1 Basic definitions

**Graph** Let  $\mathcal{G} = (\mathcal{V}, \mathcal{E})$  be an undirected graph consisting of a finite set of nodes  $\mathcal{V}$  and a set of edges  $\mathcal{E} \subseteq \binom{\mathcal{V}}{2}$ . The set  $\mathcal{E}$  will also be called a *neighborhood structure* of  $\mathcal{V}$ . For convenience, we will typically use lower case letters  $u$  and  $v$  for nodes of the graph, and write  $uv$  to denote an edge  $\{u, v\} \in \mathcal{E}$  connecting  $u$  and  $v$ . Since the graph is undirected,  $uv$  and  $vu$  denote the same edge. The notation  $\mathcal{N}_b(u)$  will be used for the set of nodes  $\{v \mid uv \in \mathcal{E}\}$  connected to the node  $u$ .

The graph  $\mathcal{G}$  is considered as a model of the considered target object, where the nodes represent the elementary object parts and edges stand for mutually dependencies between them.

In Examples 1.1 and 1.2 the graph  $\mathcal{G}$  may have the grid structure of the underlying two-dimensional pixel array. In Example 1.3 cells of one image frame are neighbors, since their labels depend on each other.

**Labels and unary costs** A finite set of labels  $\mathcal{Y}_u$  is associated with each node  $u \in \mathcal{V}$ . Our preference for each label is expressed by the *unary cost function*  $\theta_u: \mathcal{Y}_u \rightarrow \mathbb{R}$ , which is defined for each node  $u \in \mathcal{V}$ . The value  $\theta_u(s)$  determines the *cost*, which we pay for assigning label  $s \in \mathcal{Y}_u$  to the node  $u$ . Sometimes we will use very high costs to implicitly forbid

certain labels oder label pairs. The notation  $\infty$  will be used to denote such high costs.

Unary costs are usually defined by what is known from observation. In Example 1.1, typically, the color distribution in the vicinity of a given pixel defines the cost of each possible label. The difference between color distributions from two or more images of the same scene taken from different viewpoints determines the unary costs in Example 1.2. Unary costs are often called the “data term” to emphasize that they depend on the input data or observation.

**Dependence and pairwise costs** Dependencies between labels assigned to different graph nodes are modeled with *pairwise cost functions*  $\theta_{uv}: \mathcal{Y}_u \times \mathcal{Y}_v \rightarrow \mathbb{R}$ , which are defined for each edge  $uv \in \mathcal{E}$  of the graph.

A simple (although not always the best) way to model the compactness assumption in Example 1.1 is to assign

$$\theta_{uv}(s, t) = \begin{cases} 0, & s = t \\ \alpha, & s \neq t \end{cases} \quad (1.1)$$

for any pair of labels  $(s, t) \in \mathcal{Y}_u \times \mathcal{Y}_v$  with some  $\alpha > 0$ . A simple way to model a smooth surface in depth reconstruction in Example 1.2 is to assign

$$\theta_{uv}(s, t) = |s - t|, \quad (1.2)$$

to penalize large differences between depth in the neighboring nodes.

In the cell tracking example the pairwise costs should forbid the same labels to be assigned to neighboring nodes when no cell division happens:

$$\theta_{uv}(s, t) = \begin{cases} 0, & s \neq t \\ \infty, & s = t. \end{cases} \quad (1.3)$$

This disallows that cells  $u$  and  $v$  “glue” to the same “parent” cell  $s = t$ . In case of cell division, this pairwise cost function can be extended in a natural way to disallow intersection of cell descendants.

These examples show that pairwise costs often incorporate the prior information about a considered object, therefore, they are often

collectively referred to as the *prior*. However, this is not always the case. For instance, much better segmentation results can be obtained if the parameter  $\alpha$  in (1.1) depends on the color distribution of the input image, i.e. on  $uv$ .

Costs and cost functions are also called *potentials* and *potential functions*. We prefer the term *cost* since it is more widely used in general optimization literature.

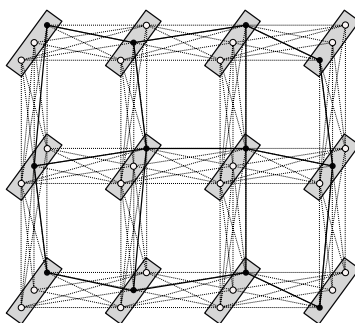
Since unary and pairwise costs are functions of discrete variables, they can be seen as vectors. Therefore we can treat the unary cost function  $\theta_u$  as a *unary cost* vector  $(\theta_u(s), s \in \mathcal{Y}_u)$ . Similar reasoning holds also for each pairwise cost function, which can be considered as a *pairwise cost* vector  $\theta_{uv} = (\theta_{uv}(s, t), (s, t) \in \mathcal{Y}_u \times \mathcal{Y}_v)$ . Unless we use the word *vector* or *function*, the context will determine whether we refer to a vector or a function  $\theta_u$  (or  $\theta_{uv}$ ). All unary vectors stacked together form the vector of all unary costs  $\theta_{\mathcal{V}} = (\theta_u, u \in \mathcal{V})$ . The vector  $\theta_{\mathcal{E}}$  of all pairwise costs is defined similarly as  $(\theta_{uv}, uv \in \mathcal{E})$ . Stacking together the latter two results in a long *cost vector*  $\theta = (\theta_{\mathcal{V}}, \theta_{\mathcal{E}})$  with dimension  $\mathcal{I} := \sum_{u \in \mathcal{V}} |\mathcal{Y}_u| + \sum_{uv \in \mathcal{E}} |\mathcal{Y}_{uv}|$ .

**Labeling** In the following, we will often use the notation  $\mathcal{Y}_{\mathcal{A}}$  for all possible label assignments to a subset of nodes  $\mathcal{A} \subseteq \mathcal{V}$ . Formally,  $\mathcal{Y}_{\mathcal{A}}$  stands for the Cartesian product  $\prod_{u \in \mathcal{A}} \mathcal{Y}_u$ . In particular,  $\mathcal{Y}_{uv}$  denotes  $\mathcal{Y}_u \times \mathcal{Y}_v$  and is the set of all possible pairs of labels in nodes  $u$  and  $v$ . A vector  $y \in \mathcal{Y}_{\mathcal{V}}$  of labels assigned to *all* nodes of the graph is called *labeling*. We will refer to coordinates of this vector with the node index, i.e.  $y_u$  stands for the label assigned to the node  $u$ . One may also speak about *partial labelings*, if only a subset  $\mathcal{A}$  of the nodes is labeled.

**Definition 1.4** (Graphical model). The triple  $(\mathcal{G}, \mathcal{Y}_{\mathcal{V}}, \theta)$  consisting of a graph  $\mathcal{G}$ , discrete space of all labelings  $\mathcal{Y}_{\mathcal{V}}$  and a corresponding cost vector  $\theta$ , is called a *graphical model*.

**Definition 1.5** (Energy minimization problem). The problem

$$y^* = \arg \min_{y \in \mathcal{Y}_{\mathcal{V}}} \left[ E(y; \theta) := \sum_{u \in \mathcal{V}} \theta_u(y_u) + \sum_{uv \in \mathcal{E}} \theta_{uv}(y_u, y_v) \right] \quad (1.4)$$



**Figure 1.2:** Labeling of the graphical model from Figure 1.1. Selected labels are marked as black circles and connected with solid lines. Each black circle corresponds to a unary cost and each solid line to a pairwise cost in the sum in the energy minimization problem (1.4).

of finding a labeling  $y^*$  with minimal total cost will be called *energy minimization* or *maximum a posteriori (MAP) inference* problem for the graphical model  $(\mathcal{G}, \mathcal{Y}_{\mathcal{V}}, \theta)$ .

For the sake of notation we will sometimes use the short form of (1.4)

$$y^* = \arg \min_{y \in \mathcal{Y}_{\mathcal{V}}} \left[ E(y; \theta) := \sum_{w \in \mathcal{V} \cup \mathcal{E}} \theta_w(y_w) \right] \quad (1.5)$$

with  $y_w$  being equal to  $y_u$ , if  $w$  corresponds to a node, i.e.  $w = u \in \mathcal{V}$ , and  $y_{uv}$ , if  $w$  corresponds to an edge, i.e.  $w = uv \in \mathcal{E}$ .

Problems equivalent or very closely related to (1.4) have also other names depending on the corresponding community they are studied in: *maximum likelihood explanation (MLE) inference* (machine learning, natural language processing community), *weighted/valued/partial constraint satisfaction problem* (constraint satisfaction community).

## 1.2 Probabilistic interpretation

The name *MAP-inference* stems from the probabilistic interpretation of the problem (1.4). With the energy  $E(y; \theta)$  one typically associates the exponential probability distribution

$$p(y) = \frac{1}{Z(\theta)} \exp(-E(y; \theta)), \quad (1.6)$$

where the normalizer  $Z(\theta)$  is known as *partition function*.

According to the distribution (1.6), problem (1.4) is equivalent to finding *the most probable* labeling  $y$ , i.e. the one maximizing  $p(y)$ . Since  $E$  has the separable form (1.5) the expression (1.6) takes the form of the *Gibbs distribution*

$$p(y) = \frac{1}{Z(\theta)} \prod_{w \in \mathcal{V} \cup \mathcal{E}} \Theta_w(y_w) \quad (1.7)$$

with  $\Theta_w = \exp(-\theta_w)$ . This explains the also frequently used name “factors” for the cost functions and their exponents  $\Theta_w$ .

The probabilistic interpretation (1.6) gives rise to several other *probabilistic inference* problems motivated by Bayesian statistical decision making theory. One computational problem, often referred to as *marginalization inference*, consists of computing *marginal* distributions

$$\hat{p}_u(s) := \sum_{y \in \mathcal{Y}_{\mathcal{V}} : y_u = s} p(y) \quad (1.8)$$

for each node  $u$  and label  $s$  of a graphical model. These kinds of problems, although closely related to MAP-inference, are beyond the scope of this monograph.

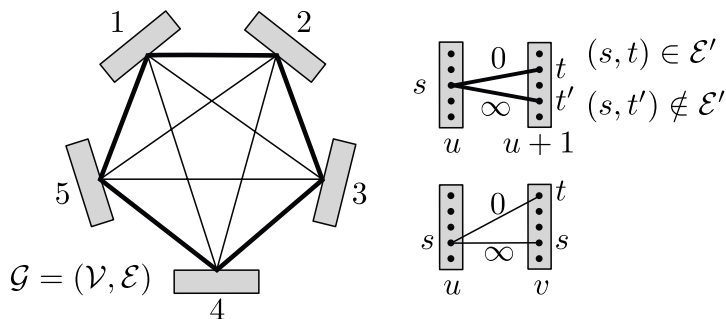
### 1.3 Combinatorial complexity of MAP-inference

The number of possible labelings  $y$  in (1.4) grows exponentially with the cardinality of  $\mathcal{V}$ , as it is equals  $\prod_{v \in \mathcal{V}} |\mathcal{Y}_v|$ . It results in  $L^{|\mathcal{V}|}$  in case all nodes have the same number of labels  $|\mathcal{Y}_u| = L, \forall u \in \mathcal{V}$ .

However, an exponentially large set of solutions is not sufficient for polynomial  $\mathcal{NP}$ -hardness of a problem. For example, the shortest path between two nodes in a directed graph with positive edge weights has an exponentially large set of solutions, but is polynomially solvable by Dijkstra’s algorithm.

Below we show that the MAP-inference (1.4) is, indeed,  $\mathcal{NP}$ -hard. To do so, it is sufficient to show that some  $\mathcal{NP}$ -complete decision problem is polynomially reducible to MAP-inference.

In the following construction we will show that the Hamiltonian cycle problem reduces to MAP-inference in polynomial time.



**Figure 1.3:** Illustration of the reduction of the Hamiltonian cycle problem to MAP-inference for a graph with 5 nodes. Edges of the MAP-inference graph  $\mathcal{G}$  are divided into two groups: between nodes  $u$  and  $u + 1$  (bold edges) and all others (thin edges). The corresponding pairwise costs are illustrated on the right (see also main text).

**Definition 1.6** (Hamiltonian cycle). A *Hamiltonian cycle* in a graph  $\mathcal{G}$  is a cycle which visits each node exactly once.

The problem of deciding whether a given directed graph has a Hamiltonian cycle is known to be  $\mathcal{NP}$ -complete. To show  $\mathcal{NP}$ -hardness of MAP-inference, it is sufficient to reduce the Hamiltonian cycle problem to it.

Let  $\mathcal{G}' = (\mathcal{V}', \mathcal{E}')$  be the graph for which one should solve the Hamiltonian cycle problem. Let us construct the following graphical model (see Figure 1.3): For the graph  $\mathcal{G} = (\mathcal{V}, \mathcal{E})$  it holds that  $\mathcal{V} = \mathcal{V}'$  and  $\mathcal{E} = \binom{\mathcal{V}}{2}$ . In other words, graph  $\mathcal{G}$  contains the same nodes as graph  $\mathcal{G}'$  and is fully connected. Moreover, we will order all nodes of the graph  $\mathcal{G}$ , i.e.  $\mathcal{V} = \{1, 2, \dots, |\mathcal{V}|\}$ . This order is the order of nodes in the Hamiltonian cycle we are searching for. We will assume the operation  $u + 1$  to be defined modulo  $|\mathcal{V}|$ , i.e.  $u + 1$  defines the next element of the Hamiltonian cycle. In other words, if  $u < |\mathcal{V}|$  then  $u + 1$  is the next natural number after  $u$  and for  $u = |\mathcal{V}|$  the element  $u + 1$  is equal to 1.

The set of labels  $\mathcal{Y}_u := \mathcal{V}'$  is the same for each node  $u \in \mathcal{V}$ . Its elements index nodes of the graph  $\mathcal{G}'$ . A label  $s$  assigned to a node  $u \in \mathcal{V}$  encodes that the  $u$ -th node in the Hamiltonian cycle corresponds to the node  $s$  of the graph  $\mathcal{G}'$ .

Unary costs are equal to 0. Pairwise costs are split into two groups. For a pair of nodes  $\{u, u + 1\} \in \mathcal{E}$  the cost reads

$$\theta_{u,u+1}(s, t) = \begin{cases} 0, & (s, t) \in \mathcal{E}' \\ \infty, & (s, t) \notin \mathcal{E}' \end{cases} \quad (1.9)$$

It guarantees that two neighboring nodes of the Hamiltonian cycle are connected by an edge in the graph  $\mathcal{G}'$ .

To guarantee that no node is included twice in the Hamiltonian cycle, we set up other pairwise costs for  $v \neq u + 1$  and  $u \neq v + 1$  as follows:

$$\theta_{uv}(s, t) = \begin{cases} 0, & s \neq t \\ \infty, & s = t \end{cases} \quad (1.10)$$

Such type of pairwise costs is sometimes called the *uniqueness constraints*, since these costs enforce that each label is selected at most ones.

Let  $y$  be some labeling of the graphical model  $\mathcal{G}$  such that  $E(y, \theta) < \infty$ . Then the sequence  $(y_1, y_2, \dots, y_{|\mathcal{V}|})$  is the Hamiltonian cycle by construction: there is an edge between  $y_u$  and  $y_{u+1}$  in  $\mathcal{G}'$ , and the set  $\{y_1, y_2, \dots, y_{|\mathcal{V}|}\}$  is exactly the set  $\mathcal{V}'$ .

All labelings have either value 0 or  $\infty$ . Therefore, the solution of the MAP-inference problem answers the question whether there is a labeling  $y$  such that  $E(y, \theta) < \infty$ , and, therefore, whether there is a Hamiltonian cycle in the graph  $\mathcal{G}'$ .

Note that the same reduction of the Hamiltonian cycle problem to the MAP-inference could have also be done without using the infinite costs. Instead, any positive finite cost (e.g. 1) could be used in place of infinities. In this case the solution of the MAP-inference problem answers the question whether there is a labeling  $y$  such that  $E(y, \theta) = 0$ , which is equivalent to the existence of a Hamiltonian cycle in the graph  $\mathcal{G}'$ .

## 1.4 Bibliography and further reading

For further examples of applications of graphical models in computer vision and image processing we refer to the collection [10]. Books [77, 47] can be recommended to learn more about the probabilistic view on



graphical models. The monograph [135] concentrates on the exponential family (1.6) of distributions and its relation to graphical models.

A classical source to learn about the computational complexity of combinatorial problems is [27], a modern exposition is given in [5]. The most recent and comprehensive analysis of complexity of the MAP-inference problem is provided in [70].

The text books [25] and [103] can be recommended to learn about Bayesian decision theory.

The reduction of the Hamiltonian cycle problem to MAP-inference is reproduced from the lectures on structural pattern recognition given by Prof. Michail Schlesinger at National Technical University of Ukraine “Igor Sikorsky Kyiv Polytechnic Institute” where the author studied mathematics and computer science.

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