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# Hardware Platform Security for Mobile Devices

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# Hardware Platform Security for Mobile Devices

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## ABSTRACT

Today, personal mobile devices like smartphones and tablets are ubiquitous. People use mobile devices for fun, for work, and for organizing and managing their lives, including their finances. This became possible because over the last two decades, mobile phones evolved from closed platforms intended for voice calls and messaging to open platforms whose functionality can be extended in myriad ways by third party developers. Such wide-ranging scope of use also means widely different security and privacy requirements for those uses. The mobile device ecosystem involved multiple different stakeholders such as mobile network operators, regulators, enterprise information technology administrators, and of course ordinary users. So, as mobile platforms became gradually open, *platform security* mechanisms were incorporated into their architectures so that the security and privacy requirements of all stakeholders could be met. Platform security mechanisms help to isolate applications from one another, protect persistent data and other on-device resources (like access to location or peripherals), and

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help strengthen software against a variety of attack vectors. All major mobile platforms incorporate comprehensive software and hardware platform security architectures, including mechanisms like *trusted execution environments (TEEs)*.

Over the past decade, mobile devices have been undergoing convergences in multiple dimensions. The distinction between “mobile” and “fixed” devices has blurred. Similar security mechanisms and concepts are being used across different platforms, leading to similar security architectures. Hardware enablers used to support platform security have gradually matured. At the same time, there have also been novel types of attacks, ranging from software attacks like return- and data-oriented programming to hardware attacks like side channels that exploit micro-architectural phenomena. It is no longer tenable to assume that the current hardware security mechanisms underpinning mobile platform security are inviolable.

The time is therefore right to take a new look at mobile platform security, which brings us to this monograph. We focus on *hardware platform security*. The monograph is divided into four parts: we begin by looking at the *why* and *how* of mobile platform security, followed by a discussion on *vulnerabilities and attacks*; we conclude by *looking forward* discussing emerging research that explores ways of dealing with hardware compromise, and building blocks for the next generation of hardware platform security.

Our intent is to provide a broad overview of the current state of practice and a glimpse of possible research directions that can be of use to practitioners, decision makers, and researchers.

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## **Part I**

# **Mobile Platform Security: Why?**

# 1

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## Introduction

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Today, mobile devices such as smartphones and tablets are very widely deployed. All modern mobile device platforms incorporate sophisticated software and hardware *platform security* mechanisms. To understand how this came to be, we need to start in the late 1990s.

### 1.1 What motivated mobile platform security?

The mobile phone revolution was well under way by the mid 1990s. Initially, mobile phones were simple embedded devices with fixed functionality: voice calls and text messages. Early on, the mobile phone industry recognized the power of billions of people having general-purpose computing devices in their hands. Personal digital assistants were already available and demonstrated the range of uses for portable, personal general-purpose computing devices. Therefore, already by the mid 1990s, the industry was working towards *opening up* mobile phone platforms so that users gain the ability to *extend* their functionality, by installing third-party software modules. This development directly led to today's smartphones and tablets with their "app" ecosystems.

The industry saw the potential for new types of applications like mobile payments, public transport ticketing, and digital media consump-

tion. But it also realized that for these applications to succeed, open mobile devices needed additional mechanisms to safeguard the security and privacy requirements of these novel, and potentially high-value, applications.

Furthermore, the mobile phone ecosystem already had well-established stakeholders. They were sensitive to the security and privacy concern that could arise in the transition from closed fixed-function devices to open platforms. Existing commodity general-purpose computing platforms at the time, like those for personal computers, did not incorporate the platform security mechanisms necessary to address these concerns. Consequently, they wanted to mediate this transition so that their own interests were safeguarded. This, too, drove the development of new platform security mechanisms. Mobile platform security architectures emerged because of the need to address these stakeholder concerns as mobile devices opened up (Matala *et al.*, 2019).

## 1.2 Stakeholders

An important class of stakeholders are mobile network operators (MNOs) (also known as “carriers”) who are motivated by business interests. An example of a business interest of MNOs is the need to strongly authenticate their subscribers. This need led to the introduction of subscriber identity modules (SIMs) (discussed further in Section 2.2) right from the beginning. Another example of a business interest of MNOs is the need for robust technical mechanisms to support the *subsidy-lock business model* where a MNO gives a mobile phone to a subscriber for free or below cost, in return for a commitment to maintain the subscription for a specified period of time. The requirement to technically enforce subsidy locks translated into each device having an unforgeable unique identifier and the ability to run subsidy-lock enforcement software in a manner that cannot be bypassed.

Another, equally important, class of stakeholders are *regulators* who safeguard the public good. An example of a regulatory need is to ensure that radio-frequency transmission parameters, which are typically calibrated for each device at the time of manufacture, cannot be tampered with. This need can be met with secure (integrity-protected) storage for storing these parameters.

A third class of stakeholders are end-users. They were used to mobile phones that were reliable and trustworthy. They expected the same degree of reliability and trustworthiness to be maintained, even as mobile phone platforms were opening up.

There are other stakeholders in the ecosystem, like enterprise administrators, and of course the mobile phone manufacturers —also known as original equipment manufacturers (OEMs)—and operating system (OS) vendors themselves. To see what kinds of mechanisms are needed to protect the interest of different stakeholders, it is necessary to understand the threat models from the perspectives of these stakeholders.

### 1.3 Threat models

A threat model involves characterizing the adversary in terms of its capabilities, and the assets that need to be protected from these adversaries. For example a *software adversary* is assumed to be capable of influencing one or more software modules on the victim device. The adversary's control may be limited to a single application (software in user space) or can extend to privileged software like the OS itself. In contrast, a *hardware adversary* can directly interact with, and possibly manipulate, the hardware components on the victim device.

Rather than presenting an exhaustive treatment of all possible threat models, we will illustrate the concept with three informal examples.

First, consider the threat of a user's address book being exfiltrated from the device by a malicious third-party application that the user happened to install. We are concerned with a software adversary (the third-party developer) and the asset that needs protection is the address book. Standard hardware support (for memory management and process isolation) combined with a good OS security architecture (providing access-controlled persistent storage for each application) would be sufficient to provide the required protection. In Section 3 we will discuss OS security architectures.

Next, consider the same setting as above, but with a different asset: credentials for accessing financial transactions like online banking. While we are still concerned with a software adversary, the value of the asset is significantly higher, and its compromise can result in substantial

losses. Consequently, relying only on OS security is not reasonable because an OS is a complex software component with a large threat surface for the attacker to exploit. Additional hardware support for protecting high-value assets is justifiable. Hardware-assisted trusted execution environments (TEEs) allow small pieces of trusted software on a general-purpose computing device to be *isolated* from the rest of the software on the same device, including the OS and other applications. Today TEEs are ubiquitous. Nearly every smartphone or tablet is likely to have a processor with TEE capabilities. Many personal computers are also equipped with TEEs. The ubiquity of TEEs is not a recent phenomenon (Ekberg *et al.*, 2014): hardware-assisted TEEs started to appear in mobile phones from the early 2000s. For a technology that is so widely deployed, for so long, the origins and trajectory of TEE technologies are poorly understood. Our primary focus in this monograph is to explore hardware platform security for mobile devices, with a particular emphasis on TEEs.

Finally, consider the case of technical mechanisms for subsidy-lock enforcement. The adversary in this case is the user of the device who has physical access to the device. The asset that the adversary wants to compromise is the binding between the mobile device hardware and the MNO (so that a successful attack will result in breaking the binding, allowing the adversary to use the device with a different MNO subscription). OS security alone is not sufficient. Since we now deal with a potential hardware adversary, we must use hardware-security mechanisms that can withstand physical attack.

#### 1.4 Chains of trust

In a given scenario, the party relying on the protection mechanism trusts the software and hardware components used to realize the mechanism. A *chain of trust* refers to the process of building up this trust, starting from one or more *roots of trust*. In the first example above, while OS security is sufficient, the relying party, the user, needs to trust that the correct OS is running on the device. *Platform integrity* (Section 4) makes it possible to build up this trust. Higher level platform security mechanisms like OS security rely on underlying building blocks like



platform integrity (Section 4), hardware-assisted isolation (Section 5), and cryptographic primitives realized in hardware (Section 6).

An important feature of hardware platform security mechanisms is allowing remote relying parties to build up trust in a device. In the second example above, a bank may need to convince itself that the user is accessing her bank account from a secure device before allowing access. This feature is called *remote attestation*, which is widely supported by modern TEEs. In Section 5 we will discuss the chains of trust involved in remote attestation in modern TEEs.

We begin with an overview of the history of mobile hardware platform security mechanisms (Section 2), and provide an overview of OS security (Section 3) to understand how an OS can make use of these mechanisms. We will explore the nuts and bolts of how platform security is implemented in today's devices, focusing on hardware platform security (Part II), and discuss attacks against hardware platform security mechanisms (Part III). We will conclude with a brief foray into a future outlook for hardware platform security (Part IV).

### Notes on the scope of this monograph

The focus of this monograph is on hardware platform security in mobile devices. We do cover OS security in Section 3, but from the perspective of motivating hardware platform security. Mobile device platforms also incorporate sophisticated *software platform security* mechanisms. We refer readers interested in this topic to books dedicated to the topic such as Asokan *et al.* (2014). We also do not cover specific high-level attacks such as *jail-breaking* (removing manufacturer-imposed restrictions on what software can be installed on a mobile device) or *rooting* (obtaining the privileges of the maximally privileged “root” user on Unix-based mobile OSs). However, the basic attacks we describe in Part III can be, and often are, used as stepping stones for these high-level attacks.

# 2

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## Historical Overview

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The requirements we saw in Section 1 led to mobile device and platform vendors developing and deploying software and hardware platform security architectures. Nokia Radio Application Processors are believed to be the first trusted execution environments (TEEs) deployed at a large scale (Matala *et al.*, 2019). These were followed shortly by Texas Instruments' M-Shield™ (Sundaresan, 2003) and subsequently by ARM's TrustZone™ (Alves and Felton, 2004) which represents the overwhelming share of deployed mobile TEEs today.

In the non-mobile setting, hardware security modules (HSMs) used in the financial sector (starting with IBM's CryptoCard<sup>1</sup>) are an early example of a TEE. Trusted Computing Group's Trusted Platform Modules (TPMs) (Arthur and Challener, 2015) are widely deployed in personal computers, where they are used for boot integrity and disk encryption, but they have not found common use in the mobile space. Recently, Intel's Software Guard Extensions (SGX) (McKeen *et al.*, 2016) has become the most widely studied TEE architecture, thanks to the easy availability of both the software and hardware.<sup>2</sup> SGX is

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<sup>1</sup><https://www.ibm.com/security/cryptocards/>

<sup>2</sup><https://software.intel.com/en-us/sgx>

primarily deployed in cloud settings to enable confidential computing use cases (Alibaba, 2020; McReynolds, 2021). Desktop use cases for SGX include Blu-ray digital rights management (DRM) (Toulas, 2021)

## **2.1 Hardware security modules**

Early examples of the inclusion of a dedicated security co-processor were motivated by the need to perform sensitive cryptographic operations isolated from other computations in systems handling financial transactions. Transaction processing for Europay, Mastercard and Visa (EMV) payment cards use HSMs as the primary security device for key management (Cryptomathic, 2017). An HSM is a discrete computing device usually encapsulated in tamper-evident coating. HSMs in backend systems typically include specialized cryptographic hardware accelerators to enable high throughput because they need to process transactions in real-time. An HSM can be realized as either a stand-alone peripheral device or as an extension board connected directly to the internal bus of the host computer. The operational keys are generated in the cryptographic co-processor within the HSM and are then saved either in a keystore file or in application memory, encrypted under the master key of that co-processor. Any HSM with an identical master key can use those keys.

The first commercially available civilian HSMs were deployed already in the 1970s, originally for IBM mainframes. The IBM 3845 and 3846 data encryption devices (IBM, 1977) allowed exported encryption keys to be encrypted using the recently standardized DES algorithm. These early HSMs included secure key entry devices (cards and PIN pads) for master key loading, random number generation capabilities for seeding, and persistent storage for key materials. They were instrumental in securing early electronic banking, such as automatic teller machines (ATMs).

### **2.1.1 HSMs in radio communication**

HSMs are also extensively deployed for modern military software-defined radio (SDR) communication. SDR refers to wireless communications

where the transmitter and receiver mixing, filtering, amplification, modulation/demodulation etc. occur in software instead of in conventional radio electronics. With SDR, software-based transmission algorithms can be downloaded and adapted over the lifecycle of the hardware. While analog military radio equipment include dedicated cryptographic chips for (proprietary) ciphers that are required for communication with compatible equipment, SDR equipment have to support a large number of cryptographic schemes, including legacy protocols and algorithms. Consequently military SDR equipment, such as the U.S. Joint Tactical Radio System (JTRS), employ embeddable HSMs specifically designed for communication security. The Advanced INFOSEC Machine (AIM) (General Dynamics Mission Systems, 2015a) is one such programmable, embeddable cryptographic unit developed by Motorola in the late 1990s. It consists of a hardware platform with three independent cryptographic processors, one for key management and two programmable processors for traffic encryption/decryption. The key management cryptographic engine (KMCE) is based on a 32-bit reduced instruction set computer (RISC) processor and includes a math co-processor designed for public key algorithm processing. The KMCE runs a read-only memory (ROM)-based Secure Operating System (SOS). The SOS provides a multi-security level, multi-tasking environment for the cryptographic applications which allowed the functionality of the AIM to be extended by software. The chip contains the necessary building blocks to implement encryption algorithms such as DES, and the classified SAVILLE and BATON cryptographic algorithms used by U.S. and NATO. Its successor, AIM II (General Dynamics Mission Systems, 2015b) is specifically designed for JTRS. Around the same time, a similar crypto-chip, called the General Crypto Device (GCD) (Lange, 1997), was developed in Europe by Dutch electronics giant Philips.

The use of HSMs such as AIM and GCD are early examples of the use of TEEs in telecommunications. The sensitivity of military communication justified the inclusion of dedicated components for security into end devices. However, for civilian telecommunication devices, the widespread use of TEE-technology only occurred when two conditions were met: 1) economic incentives emerged to justify requiring strong, hardware-based security, and 2) low-cost technological solutions that met those requirement were developed.

## 2.2 SIMs, mobile handsets, and smart cards

During the early 1990s, civilian wireless communication systems also began to employ hardware-assisted security. Mobile network operators (MNOs) required a reliable way of preventing illicit use of a subscriber identity for making phone calls from mobile phones. For this purpose, the subscriber identity module (SIM) card (GSMA, 2015) was developed by Munich smart-card maker Giesecke & Devrient, who sold the first 300 SIM cards to the Finnish MNO Radiolinja in 1991. The use of SIM cards became mandatory in the Global System for Mobile Communications (GSM) standard. Each SIM card contains an international mobile subscriber identity (IMSI) that uniquely identifies the user of the mobile network and a unique symmetric cryptographic key ( $K_i$ ) assigned to it by the MNO during SIM card personalization. The SIM ensures the integrity of the IMSI and  $K_i$ , and the confidentiality of  $K_i$ .  $K_i$  allows the MNO to authenticate the SIM card when the mobile phone connects to the network. When the mobile phone connects, it obtains the IMSI from the SIM card, and requests network access by transmitting the IMSI to the MNO. The MNO looks up the corresponding  $K_i$  of the IMSI from its subscriber database, and generates a random nonce as a challenge which is transmitted to the mobile phone. The mobile phone passes the challenge to the SIM card, which signs it, and returns the signed response, which is transmitted back to the MNO by the mobile phone. The MNO compares the signed response to the response calculated using the MNO's copy of the  $K_i$ . If they match, the authentication is successful.

Modern SIM cards are based on tamper-resistant universal integrated circuit card (UICC) technology (SIMalliance Ltd., 2013) similar to smart cards. UICC cards can host multiple software applications, typically developed using Java Card software technology (ETSI, 2012). The applications include a SIM application for GSM, and universal subscriber identity module (USIM) for UMTS (3G), Long-Term Evolution (4G), and 5G network authentication. MNOs can also provision additional value-add applications to UICC cards that they issue, such as mobile banking and phone-based money transfer. UICC application can interface with mobile phone users or initiate actions via a card

application toolkit (CAT) part of the mobile phone operating system (OS): SIM Application Toolkit (STK) for GSM systems, and USIM Application Toolkit (USAT) for later generation networks. UICCs can support an optional bearer independent protocol (BIP), which allows MNOs to deliver over-the-air (OTA) updates to UICC applications either via cell broadcasts, or short message service packets.

All UICC applications are subject to authorization by the issuer security domain (ISD), namely the MNO who issued the UICC. Consequently UICCs are effectively closed application ecosystems; it is not possible for third-party developers to leverage UICC security without co-operating with MNOs in their region. This puts add-on services operated by large MNOs into an advantageous position compared to third-party alternatives, as is the case with M-Pesa (Mbiti and Weil, 2011), a money transfer application operated by Safaricom and Vodacom, the largest mobile MNO in Kenya and Tanzania. In developing countries, such as Kenya, low-cost feature phones are still prevalent, and UICC applications is the only ubiquitous application platform available to the majority of mobile phone users. Proprietary SIM overlay technology (a.k.a. “slim SIM” or “skin SIM”) (Mondato, 2014) can enable third-party applications to operate independently of the underlying UICC.

The SIM overlay is a computer chip embedded into a thin plastic sheet that can be placed on top of a standard UICC card within a mobile phone. They were originally developed to support low-cost mobile roaming for Chinese customers traveling outside their home province. The overlay SIM acts as an independent security device, and allows additional functionality to be added to any mobile phone by attaching the overlay SIM to an MNO-issued UICC. However, an overlay SIM also has the potential to facilitate a man-in-the-middle attack by observing sensitive data such as personal identification numbers (PINs) being transmitted to the underlying UICC, or initiate, intercept and/or block mobile communications or CAT instructions (GSMA, 2014). By obtaining unauthorized access to the UICC SIM applications they could also change MNO configuration settings.

Embedded SIMs (eSIMs) (GSMA, 2015) are secure elements physically integrated into a mobile phone. eSIM chips can be directly soldered

onto the device or even embedded into the system on chip (SoC) itself. This physical integration necessitates MNO SIM or USIM profiles to be remotely provisioned. Additionally, unlike removable SIM cards, a single eSIMs may need to store multiple MNO profiles simultaneously.

### 2.3 Processor secure environments

Towards the late 1990s, mobile phones were transitioning from closed systems to open application platforms, for which third-party applications could be developed using the Java programming language. While not yet true smartphones, the feature phones of the time were gradually starting to resemble small, general-purpose computers. This brought with it new business opportunities, but also new challenges for device security; regulators and MNOs needed to ensure the protection of certain pieces of information after the mobile phone had left the manufacturing line. In particular, regulators required that the device identity, the international mobile equipment identifier (IMEI), remain unchanged in order to act as a theft deterrent. IMEIs of stolen mobile phones are blacklisted by network operators, thereby reducing the economic value of stolen mobile phones and deterring theft (GSMA, 2019).

Similarly, radio frequency parameters, which could affect the quality of service of other mobile phones in the area, or the safety of the user, should also remain unchanged. MNOs, who were the primary customers of large original equipment manufacturers (OEMs) such as Nokia, were concerned with ensuring that their subscribers receiving subsidized mobile devices do not break their contract terms. Consequently, they required a strong *subsidy lock* mechanism (colloquially known as SIM lock), which would tie the mobile phone to a particular MNO for the duration of the contract. Another emerging use case was DRM for digital content sold by the MNOs; initially ringtones, later games and music.

Nokia was the first to pursue a hardware-enforced processor secure environment. At the time, the security of Nokia's Digital Core Technology (DCT) generation phones was mainly based on obfuscated software solutions and protected by secrecy within the organization; even within the company, only few security professionals knew the exact design and

requirements of the DCT security architecture (Matala *et al.*, 2019). The leading market share of Nokia made it an attractive target for hackers who, (typically for a small fee) would “unlock” or “unbrand” subsidy-locked phones by either reverse engineering the valid unlock codes, or reflashing the phone with a different firmware version.

The fourth generation of DCT mobile phones included hardware components in the form of one-time-programmable memory to aid in the secure storage of sensitive device parameters. However, in the case of SIM locks, the economic motives to break device security were higher than the capabilities of the protection mechanism deployed at the time. Consequently, the revenue losses of important MNO customers resulting from SIM unlocking, increased the pressure to design a better security architecture for the upcoming generation of Nokia phones.

Within Nokia the idea of a coherent, hardware-enforced platform security originated within a team of engineers working with mobile payments and security (Matala *et al.*, 2019). Initial designs revolved around introducing a discrete security co-processor to ensure the physical isolation of the security-critical operations. However, the additional hardware chip in the bill of materials was deemed too expensive in the extremely cost-conscious organization, whose competitive advantage largely stemmed from its ability to keep manufacturing and components costs in control. Instead, Nokia engineers opted to implement a logically isolated secure processing mode within the main central processing unit (CPU). This solution was not only more cost effective in terms of component costs during manufacturing, but also functioned as common hardware platform for solutions to different use cases. This processor secure environment (Ekberg, 2013) would form the cornerstone of Nokia’s Baseband 5 (BB5) generation mobile phone security architecture.

Initial hardware designs were based on Nokia’s own radio application processors (RAPs), but from very early on Nokia collaborated with the U.S. semiconductor and Integrated Circuit (IC) manufacturer Texas Instruments (TI) with whom they had a close partnership at the time. The first BB5 mobile phone, the Nokia 6630 (codename “Charlie”) was based on TI’s Open Multimedia Applications Platform (OMAP) processors based on the ARM architecture. TI would brand the processor secure environment technology initially developed jointly with Nokia



as M-Shield (Sundaresan, 2003). It was however in Nokia's interest to ensure that it could invite bids from multiple hardware manufacturers for processors implementing a security architecture meeting Nokia's requirements. This became possible around 2003, when ARM proposed to develop system-wide hardware isolation architecture for secure execution for the ARMv6-A application processor architecture which included security extensions to the ARM SoC covering the processor, memory controllers and peripherals. ARM's design would become known as TrustZone (Alves and Felton, 2004). Integrating TrustZone in ARM processor architecture would ensure that any semiconductor manufacturer that implemented the TrustZone security extensions could supply Nokia with processor chips that met their requirements.

## 2.4 Trusted execution environments

In Section 1, we introduced the notion of TEEs – intuitively, a TEE is a computing environment on a device that a relying party trusts to a greater extent than the rest of the software running on the same device. Consider a device running a general-purpose operating system and applications, which, following standard practice, we will refer to as rich execution environment (REE) (GlobalPlatform, 2018c). For the purposes of this monograph we deem the device to have a TEE capable of running trusted code, if it has the following *capabilities*, possibly based on hardware support:

1. **Isolation:** The ability to run trusted code strongly isolated from the REE so that the REE cannot influence or learn the computations carried out by the trusted code,
2. **Secure Storage:** The ability for the trusted code to store persistent data guaranteeing its integrity and confidentiality with respect to an adversarial REE, even across reboots, and
3. **(Remote) Attestation:** The ability to convince a (possibly remote) party of the presence of the above attributes, and the characteristics of the trusted software protected by them.

This is an intentionally broad definition. It encompasses both physically distinct components—such as HSMs and TPMs—as well as processor secure environments where the isolation is logical and is enabled by extensions to the processor hardware.<sup>3</sup>

TEEs have largely evolved based on business needs, a number of commercial TEEs (Table 1) have emerged over the years. For mobile TEEs there is a framework of applicable standards, and a core set of these has reached critical mass in industry adoption. Standardization has followed in two contexts: 1) whenever and wherever common interfaces and application programming interfaces (APIs) are needed for interoperability, and 2) where common agreement for the formulation of the required security level for today’s TEEs has been required.

The main standardization organization for mobile TEEs is the GlobalPlatform (GP) consortium.<sup>4</sup> GP provides a system architecture document (GlobalPlatform, 2018c) that describes the main components of the standards set related to TEEs, and how these individual standards contribute to the overall TEE system. Ostensibly the GP TEE architecture is not tied to any particular underlying hardware mechanism for ensuring isolation, but is, in practice, heavily influenced by the ARM TrustZone security architecture. Consequently GP standards are primarily adopted by TrustZone-based TEEs. Enclave architectures (Section 5.2), such as Intel SGX, do not yet have well-defined interoperability specification. But there are on-going efforts like the Linux Foundation’s Confidential Computing Consortium which includes projects like the Open Enclave SDK<sup>5</sup> to provide a common development environment across different enclave architectures.

The GP TEE Client API (GlobalPlatform, 2010) is the common operating system interface (endpoint) to all TEE services. The specification primarily includes APIs for installing trusted applications (TAs) within the TEE, and for allowing REE applications – also known as

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<sup>3</sup>Sometimes the term TEE is used as a synonym for the particular instance that we call “processor secure environments” in Section 5. The broad definition we adapt in this monograph is consistent with the terminology used by GlobalPlatform (GlobalPlatform, 2018c).

<sup>4</sup><https://globalplatform.org/>

<sup>5</sup><https://openenclave.io/>

client applications (CAs) – to communicate with their respective TAs, defining the data interaction model and the session management for this purpose. A separate Debug API, when available, enables a TA developer to receive logs from his TA, and also some post-mortem data in the case of critical crashes.

The GP TEE Internal API (GlobalPlatform, 2021a) is the specification against which TAs are written. For the time being, it provides C-language binding. The internal API defines the transactional model of TAs in the form of a set of standardized callback functions that are called when the TA is loaded, when it is connected to initially, and when it receives an incoming command. The data formats are TA-specific, but communication follows a paradigm of shared memory, allocated by the caller and accessible by the TA, when an incoming message is received. Another aspect of the internal API is the standardized programming framework, a “libc-like” interface that provides the TA developer with memory management, secure storage, time, peripheral access and cryptographic primitives. Due to the emphasis on security, the coverage of the cryptographic functionality in the internal API is extensive, and features most contemporary algorithms for public and private key cryptography, symmetric ciphers as well as digest and signature functions. Optional extensions (standards) to the GP internal API includes interfaces to smart cards and embedded secure elements (from within the TEE) (GlobalPlatform, 2021b), APIs by which trusted user interfaces can be setup and controlled (GlobalPlatform, 2013; GlobalPlatform, 2018d), and a socket API for network endpoints (GlobalPlatform, 2017).

For remote administration of TEEs, two separate specifications exist. Both are based on the notion that security domains are established on the device in a hierarchical fashion, after which the lifecycle of a security domain can be remotely managed, and secrets (data) and TA codes can be remotely provisioned to it. The two variants are the TEE Management Framework (TMF) (GlobalPlatform, 2016), and the Open Trust Protocol (OTrP) (Pei *et al.*, 2019; GlobalPlatform, 2019b). The latter is specified both in the context of the GP consortium (GlobalPlatform, 2019b) and in the context of IETF (Pei *et al.*, 2019). Even though both protocols accomplish the same thing, TMF is better suited to off-line (or store-and-forward) provisioning, whereas OTrP is explicitly an online protocol.

Another provisioning standard, used for virtually all smart cards with application update functionality (including UICC cards) is GP's *Card Specification* standards (GlobalPlatform, 2018a). These define the card commands by which software can be provisioned to the smart cards, and how security domains, i.e., keys identifying a certain card context, are managed. The secure communication between the provisioning entity and the card, as used by the Card Specification standard, is defined in the GP Secure Channel Protocols (GlobalPlatform, 2019a).

## **Appendix**

## Commercial TEE Deployments

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Since TEE technology, and in particular TrustZone, has been deployed in large scale, a number of TEE vendors have emerged over the years. The majority of these are with proprietary implementations of the TEE software stack. Table 1 lists TEE vendors for TrustZone, TrustZone-M and the RISC-V architecture.

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<sup>1</sup>[https://globalplatform.org/wp-content/uploads/2019/07/01-CR-1.0\\_GP180004-Certificate-and-Certification-Report\\_20190712.pdf](https://globalplatform.org/wp-content/uploads/2019/07/01-CR-1.0_GP180004-Certificate-and-Certification-Report_20190712.pdf)

<sup>2</sup><https://www.commoncriteriaportal.org/files/epfiles/CC%20Huawei%20Trustee%20Software%20V2.0%20Security%20Target%202.1.pdf>

<sup>3</sup><https://www.trustonic.com/solutions/iot-security/>

<sup>4</sup><https://www.op-tee.org/>

<sup>5</sup><https://www.provenrun.com/products/provencore/>

<sup>6</sup><https://www.qualcomm.com/products/features/mobile-security-solutions>

<sup>7</sup><https://www.rockycore.cn/index.html>

<sup>8</sup><https://optimumdesk.com/it-solutions/data-loss-prevention-privacy>

<sup>9</sup><https://www.sierraware.com/open-source-ARM-TrustZone.html>

<sup>10</sup><https://www.trustkernel.com/en/products/tee/t6.html>

<sup>11</sup><https://developer.samsung.com/teegris>

<sup>12</sup><https://source.android.com/security/trusty>

<sup>13</sup>[https://www.taf.net.cn/Tee\\_detail.aspx?\\_ID=2349283b-6311-4617-862c-112234544354](https://www.taf.net.cn/Tee_detail.aspx?_ID=2349283b-6311-4617-862c-112234544354)

<sup>14</sup><https://globalplatform.org/certified-products/watchtrust-2-1-1-on-sc9860-2/>

<sup>15</sup>[https://globalplatform.org/wp-content/uploads/2019/09/GP-SE-2019\\_02-CR-1.0\\_GP190006-Certificate-and-Certification-Report\\_20190909.pdf](https://globalplatform.org/wp-content/uploads/2019/09/GP-SE-2019_02-CR-1.0_GP190006-Certificate-and-Certification-Report_20190909.pdf)

<sup>16</sup><https://www.trustonic.com/technical-articles/kinibi-m/>

<sup>17</sup><https://www.st.com/en/embedded-software/provencore-m.html>

<sup>18</sup><https://hex-five.com/first-secure-iot-stack-riscv/>

**Table 1:** Commercial TEE implementations. The GlobalPlatform (GP) Certification column indicates TEEs that have received TEE Initial Configuration v1.1 functional certification and/or TEE Security certification according to <https://globalplatform.org/certified-products/>.

TEE	Vendor	GP Certification Functional Security	License	Note
<b>TrustZone</b>				
Clond Link TEE <sup>1</sup>	Pingtong Semiconductor	✓	proprietary	
iTrustee <sup>2</sup>	Huawei	✓	proprietary	Formerly Secure Core
Kinibi <sup>3</sup>	Trustonic	✓	proprietary	Formerly Mobicore and <t-base
OP-TEE <sup>4</sup>	Linaro		BSD 2-Clause	
ProvenCore <sup>5</sup>	Prove & Run		proprietary	
Qualcomm TEE <sup>6</sup>	Qualcomm		proprietary	Formerly QSEE
Rocky Core <sup>7</sup>	Suzhou Rong Card Intelligent Technology	✓	proprietary	
SecuriTEE <sup>8</sup>	Solacia		proprietary	
SierraTEE <sup>9</sup>	SierraWare		proprietary	
T6 <sup>10</sup>	TrustKernel		GNU GPL	
TEEGRIS <sup>11</sup>	Samsung	✓	proprietary	GP certification for Mediatek MT6737T
TLK NVIDIA Corporation, 2015	Nvidia		MIT	
Trusty <sup>12</sup>	Google		MIT	
TURBOTEE <sup>13</sup>	Eastcompeace Technologies	✓	proprietary	
Ymos TEE	Taobao Software	✓	proprietary	
WATCHTRUST <sup>14</sup>	Watchdata	✓	proprietary	
Upteq NFC422 v1.0 <sup>15</sup>	Gemalto (Thales Group)	✓	proprietary	
<b>TrustZone-M</b>				
Kinibi-M <sup>16</sup>	Trustonic	N/A	proprietary	
ProvenCore-M <sup>17</sup>	Prove & Run	N/A	proprietary	
<b>RISC-V</b>				
MultiZone <sup>18</sup>	Hex Five Security	N/A	MIT	

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